

Exercise 1

(10+2+2=14 Marks)

A robot wants to play a guessing game with numbers against a computer. The goal for the robot is to guess a given number between two numbers as quickly as possible. The algorithm should display the guessed number along with the number of guesses.

- a) You are asked to write an algorithm that given a number N which is initially hidden for the robot, it finds this number in the **least** number of guesses. The algorithm takes as input the hidden number and the range.

For example: For a given number 14 and a range between 1 and 50 the algorithm should display the following

```
Please guess a number between 1 and 50
The robot guesses 25
The computer says my number is lower than your guessed one
The robot guesses 12
The computer says my number is higher than your guessed one
The robot guesses 18
The computer says my number is lower than your guessed one
The robot guesses 15
The computer says my number is lower than your guessed one
The robot guesses 14
The robot found the number in 5 guesses
```

Solution:

```
get CN
get lo
get hi
print "Please guess a number between " + lo + "and " + hi
set GN to INT((lo + hi)/2)
print "The robot guesses " + GN
set c to 1
while(GN <> CN){
    if (CN < GN)
    then
        print "The computer says my number is lower than your guessed one"
        set hi to GN - 1
    else
        print "The computer says my number is higher than your guessed one"
        set lo to GN + 1
    endif
    set GN to INT((lo + hi)/2)
    set c to c + 1
    print "The robot guesses " + GN
}
print "The robot found the number in" + c + " guesses"
```

- b) For a given number between 1 and 100, how many guesses at most should be done according to your algorithm?

Solution:

7 guesses.

- c) For a given number between 1 and n , find out an expression that corresponds to the maximum number of guesses according to your algorithm?

Solution:

$\lceil \log_2(n) \rceil$

Exercise 2

(8+3+6+2=19 Marks)

Given the following algorithm

```

get n
get A1, ..., An
set i to 1
while(i =<= n)
{
  if(A[absolute(A[i])] > 0)
  then
    set A[absolute(A[i])] = -A[absolute(A[i])]
  else
    print absolute(A[i])
  endif
  set i to i + 1
}

```

where `absolute(X)` calculates the absolute value of `X`.

- a) What is the output of the algorithm for the following input.

1 3 2 2 1

Draw a tracing table.

Solution:

n	i	A[i]	A[absolute(A[i])]
5	1	1	1
5	2	3	2
5	3	-2	3
5	4	2	-3
5	5	1	-1

2 1

- b) What does the algorithm do for any list of length
- `n`
- consisting of elements in the range of 1 to
- `n`
- ?

Solution:

The algorithm prints all elements of the list that are duplicated. Note that if an element is duplicated k times, then this element will be printed $k - 1$ times.

c) Find the total number of executed operations. Show your workout.

```

get n
get A1, ..., An
set i to 1 -----> 1 operation --> executed once
while(i =<= n) ----->
{
  if(A[absolute(A[i])] > 0) ----->
  then
    set A[absolute(A[i])] = -A[absolute(A[i])] ---->
  else
    print absolute(A[i]) ----->
  endif
  set i to i + 1 ----->
}

```

Solution:

```

get n
get A1, ..., An
set i to 1 -----> 1 operation --> executed once
while(i =<= n) -----> 1 operation --> executed (n + 1) times
{
  if(A[absolute(A[i])] > 0) -----> 1 operation --> executed (n) times
  then
    set A[absolute(A[i])] = -A[absolute(A[i])] ----> 1 operation --> executed once
  else
    print absolute(A[i]) -----> 1 operation --> executed (n-1) times
  endif
  set i to i + 1 -----> 1 operation --> executed (n) times
}

```

Total number of operations: $1 + n + 1 + n + 1 + (n-1) + n = 4n + 2$

This holds for all lists however the we computed the total number of operations based on the case where all elements of the list have the same value.

d) Determine the order of magnitude of the algorithm.

Solution:

$O(n)$

Exercise 3

(10 Marks)

Write an algorithm that given a list of strings representing a series of coin tosses (**head** and **tail**) returns true if the list contains anywhere within it 10 consecutive occurrences of heads.

Solution:

```
get n
get A1...An
set i to 1
set count to 0

while (i <= n)
{
  if (Ai = "head")
  then
    set count to count + 1
    if(count = 10)
    then
      print "True"
      set i to n
    endif
  endif

  else
    set count to 0
  endif
  set i to i + 1
}
```

Exercise 4

(2+2+2=6 Marks)

- a) Convert the binary number 100010_2 which is in two's complement to a decimal number (base 10). Show your workout.

Solution:

The number is negative, then get its 2's complement: $011101 + 1 = 011110$, which is -30_{10}

- b) Convert the decimal number 421_{10} to a number in base 6. Show your workout.

Solution:

Division	Quotient	Remainder
421	70	1
70	11	4
11	1	5
1	0	1

1541_6

- c) Convert the hexadecimal number $C3AF15_{16}$ to a number in base 2. Show your workout.

Solution:

The number in groups of four bits: 1100 0011 1010 1111 0001 0101

Exercise 5

(3+2+3=8 Marks)

We would like to store the floating-point number -42.31 in a computer that uses 16 bits to represent real numbers. (10 for the mantissa and 6 for the exponent, both including the sign bit). Show your work as indicated below.

- a) Show the binary representation of the decimal number -42.31 .

Solution:

-101010.010

- b) Show the binary number in normalized scientific notation.

Solution:

$-.101010010 * 2^6$

- c) Show how the binary number will be stored in the 16 bits below.

Solution:

1	101010010	0	00110
Sign of mantissa 1 bit	Mantissa 9 bits	Sign of exponent 1 bit	Exponent 5 bits

Exercise 6

(6 Marks)

Assume that our computer stores decimal numbers using 6 bits. Perform the subtraction

$$(-22)_{10} - (10)_{10}$$

using 2's complement notation. Give the result of the subtraction in decimal. Show your workout, i.e. all steps performed.

Solution:

+22: 010110

-22: 101001 + 1 = 101010

+10: 001010

-10: 110101 + 1 = 110110

$(-22) + (-10) = 101010 + 110110 = 1\ 100000$

100000 get its 2's complement: $011111 + 1 = 100000_2 = -32_{10}$

Exercise 7

(4+4=8 Marks)

Given the following truth table where **A** and **B**, **C** and **D** are the input variables and **X**, **Y**, **Z**, and **W** are the output variables.

A	B	C	D	X	Y	Z	W
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	0	0	1	0	0
0	0	1	1	1	0	0	1
0	1	0	0	0	0	0	1
0	1	0	1	0	0	0	0
0	1	1	0	0	0	1	1
0	1	1	1	1	0	0	0
1	0	0	0	0	1	0	0
1	0	0	1	0	0	1	1
1	0	1	0	0	0	0	0
1	0	1	1	0	1	0	1
1	1	0	0	1	0	0	1
1	1	0	1	1	0	0	0
1	1	1	0	0	1	0	1
1	1	1	1	0	0	0	0

- a) Use the sum-of-products algorithm to find the Boolean expressions that describe the output of the truth table.

Solution:

$$X = A'B'CD + A'BCD + ABC'D' + ABC'D$$

$$Y = A'B'CD' + AB'C'D' + AB'CD + ABCD'$$

$$Z = A'BCD' + AB'C'D$$

$$W = A'B'C'D + A'B'CD + A'BC'D' + A'BCD' + AB'C'D + AB'CD + ABC'D' + ABCD'$$

- b) What is the functionality of the circuit? **Hint:** Think about **AB** and **CD** as a representation of two numbers consisting of two bits each.

Solution:

Assume that the two numbers are N_1 and N_2 with the corresponding binary representation AB and CD respectively.

The circuit computes the operation $|N_1^2 - N_2^2|$.

Exercise 8

(8+8=16 Marks)

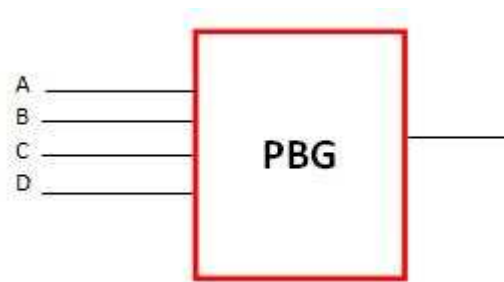
A parity bit generator is a circuit that outputs one if the number of ones in the input is odd.

- a) Your task is to construct a truth table for a 4-bit input parity generator.

Solution:

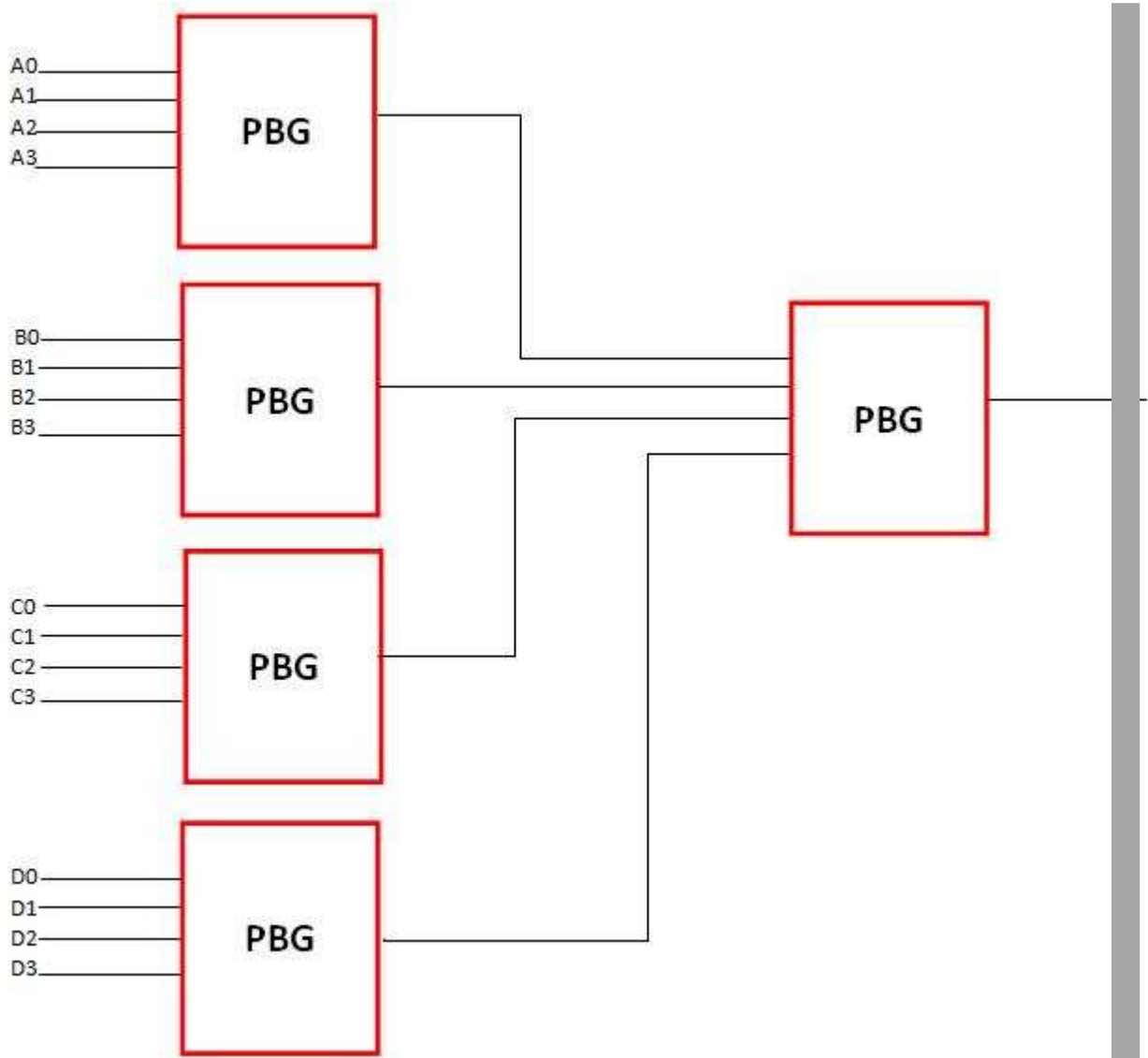
A	B	C	D	P
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
0	1	1	1	1
1	0	0	0	1
1	0	0	1	0
1	0	1	0	0
1	0	1	1	1
1	1	0	0	0
1	1	0	1	1
1	1	1	0	1
1	1	1	1	0

- b) Assume that you have already manufactured 4-bit input parity generators (PBG).



We would like to design a circuit of 16-bit input parity generator. Design a circuit using only 4-bit input parity generators as shown in the figure.

Solution:



Exercise 9

(4+8=12 Marks)

Given the following Boolean expression

$$C'D' + CB' + B'D$$

- a) Draw the circuit using and, or and not gates.
 b) Simplify the Boolean expressions using the Boolean algebra. Please mention the applied rules.

$x + 0 = x$	$x * 1 = x$	
$x + 1 = 1$	$x * 0 = 0$	
$x + x = x$	$x * x = x$	
$x + x' = 1$	$x * x' = 0$	
$(x')' = x$		
$x + y = y + x$	$xy = yx$	Commutativity
$x + (y + z) = (x + y) + z$	$x(yz) = (xy)z$	Associativity
$x(y + z) = xy + xz$	$x + yz = (x + y)(x + z)$	Distributivity
$(x + y)' = x'y'$	$(xy)' = x' + y'$	DeMorgan's Law

Hint: The circuit of the simplified expression consists of only three gates.**Solution:**

$$\begin{aligned}
 &= C'D' + B'C + B'D && (\text{Commutativity}) \\
 &= C'D' + B'(C + D) && (\text{Distributivity}) \\
 &= C'D' + B'(C + D)'' && (x'' = x) \\
 &= C'D' + B'(C'D')' && (\text{DeMorgan's Law}) \\
 &= C'D' + (C'D')'B' && (\text{Commutativity}) \\
 &= (C'D' + B')(C'D' + (C'D')') && (\text{Distributivity}) \\
 &= (C'D' + B') * 1 && (x + x' = 1) \\
 &= C'D' + B' && (x * 1 = x) \\
 &= (C + D)' + B' && (\text{DeMorgan}) \\
 &= ((C + D) * B)' && (\text{DeMorgan})
 \end{aligned}$$

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